

Design of Algorithms by Induction for Some Problems

Dr. Gur Saran Adhar

+

+

Approach: Design by Induction- Evaluating Polynomials

Problem: Given a sequence of real numbers a_n, a_{n-1}, \dots, a_0 , and a real number x , compute the value of the polynomial

$$P_n(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0.$$

Reference Udi Manber page 92

+

1

+

+

Algorithm Polynomial Evaluation(\bar{a}, x)

Input: $\bar{a} = a_0, a_1, \dots, a_n$ (coefficient of polynomial,
and x (a real number)

Output: P (the value of the polynomial at x)

begin

$$P = a_n$$

for $i = 1$ to n do

$$P = x * P + a_{n-i}$$

end.

Reference Udi Manber page 94

+

2

+

+

Approach: Design by Induction- Maximum Consecutive Subsequence

Problem: Given a sequence x_1, x_2, \dots, x_n of real numbers (not necessarily positive) find a subsequence x_i, x_{i+1}, \dots, x_j (of consecutive elements) such that the sum of the numbers in it is maximum over all subsequences of consecutive elements.

Example:

In the sequence $(2, -3, 1.5, -1, 3, -2, -3, 3)$, the maximum subsequence is $(1.5, -1, 3)$; its sum is 3.5.

If all the numbers are negative, then the maximum consecutive subsequence is empty and the sum of the empty subsequence is zero.

Reference Udi Manber page 106

+

3

+

+

Maximum Consecutive Subsequence(X, n)**Input:** x an array of size n **Output:** Global_Max

(The sum of the maximum subsequence)

begin $Global_Max = 0$ $Suffix_Max = 0$ for $i = 1$ to n do**if** $x[i] + Suffix_Max > Global_Max$ then $Suffix_Max = Suffix_Max + x[i]$ $Global_Max = Suffix_Max$ **else if** $x[i] + Suffix_Max > 0$ then $Suffix_Max = x[i] + Suffix_Max$ **else** $Suffix_Max = 0$ **end.**

reference Udi Manber page 107

+

4