

### Theorem 6.3.1

Let  $G = (V, \Sigma, P, S)$  be a regular grammar.

Define the NFA  $M = (Q, \Sigma, \delta, S, F)$  as follows:

$$i). Q = \begin{cases} V \cup \{Z\} & \text{where } Z \text{ not in } V, \text{ if } P \text{ contains a rule } A \rightarrow a \\ V & \text{otherwise} \end{cases}$$

ii).

$$\delta(A, a) = B, \text{ whenever } A \rightarrow aB \in P$$

$$\delta(A, a) = Z, \text{ whenever } A \rightarrow a \in P$$

(iii).

$$F = \begin{cases} \{A \mid A \rightarrow \lambda \in P\} \cup \{Z\} & \text{if } Z \in Q \\ \{A \mid A \rightarrow \lambda \in P\} & \text{otherwise} \end{cases}$$

Then  $L(M) = L(G)$