Transformation Algorithms

Reference Sudkamp, Chapter 4

## Algorithm 4.2.1 Construction of the NULL Set

### input

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Context Free Grammar  $G = (V, \Sigma, P, S)$ 1. $NULL = \{A \mid A \rightarrow \lambda \in P\}$ 2.repeat 2.1 PREV = NULL; 2.2 for each variable  $A \in V$  do if there is a rule  $A \rightarrow w$  and  $w \in PREV^*$   $NULL = NULL \cup \{A\}$ until NULL = PREV

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# Algorithm 4.3.1 Construction of the Set CHAIN(A) input Context Free Grammar $G = (V, \Sigma, P, S)$ 1. $CHAIN(A) = \{A\}$ 2. $PREV = \phi$ 3.repeat3.1 NEW = CHAIN(A) - PREV; 3.2 PREV = CHAIN(A)3.3 for each variable $B \in NEW$ do for each rule $B \rightarrow C$ do $CHAIN(A) = CHAIN(A) \cup \{C\}$ until CHAIN(A) == PREV

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## Algorithm 4.4.2

Construction of the Set of variables that derive Terminal Strings

#### input

Context Free Grammar  $G = (V, \Sigma, P, S)$ 1. $TERM = \{A \mid A \rightarrow w \in P\}$  with  $w \in \Sigma^*$ 2.**repeat** 

2.1 PREV = TERM; 2.2 for each variable  $A \in V$  do if there is a rule  $A \rightarrow w$ ,  $w \in \{PREV \cup \Sigma\}^*$  then  $TERM = TERM \cup \{A\}$ until PREV == TERM

## Algorithm 4.4.4 Construction of the Set of Reachable Variables

### input

Context Free Grammar  $G = (V, \Sigma, P, S)$ 1. $REACH = \{S\}$ 2. $PREV = \phi$ 3.repeat3.1 NEW = REACH - PREV; 3.2 PREV = REACH3.3 for each variable  $A \in NEW$  do for each rule  $A \rightarrow w$  do add all variables in w to REACH until REACH == PREV